

CS425

Computer Systems Architecture

Fall 2023

Static Instruction Scheduling

Techniques to reduce stalls

$CPI = \text{Ideal CPI} + \text{Structural stalls per instruction} + \text{RAW stalls per instruction} + \text{WAR stalls per instruction} + \text{WAW stalls per instruction}$

We will study two types of techniques:

Dynamic instruction scheduling	Static instruction scheduling (SW/compiler)
Scoreboard (reduce RAW stalls)	Loop Unrolling
Register Renaming (reduce WAR & WAW stalls) <ul style="list-style-type: none">• Tomasulo• Reorder buffer	SW pipelining
Branch Prediction (reduce control stalls)	Trace Scheduling

Dependencies between Instructions

- What are the sources of stalls/bubbles?
 - instructions that use the same registers
- **Parallel** instructions can execute without imposing any stalls (if we ignore structural hazards)
 - DIV.D F0, F2, F4
 - ADD.D F10, F1, F3
- **Dependencies** between instructions may lead to stalls
 - DIV.D F0, F2, F4
 - ADD.D F10, F0, F3

RAW must enter the execution stage in order
- The dependencies between instructions limit the order of execution of these instructions (impose in order execution). In the 2nd example ADD.D **must** execute after DIV.D has completed. On the other hand, parallel instructions **may** execute in the any order (out-of-order execution). In the 1st example ADD.D can execute before DIV.D.

Dependencies between Instructions

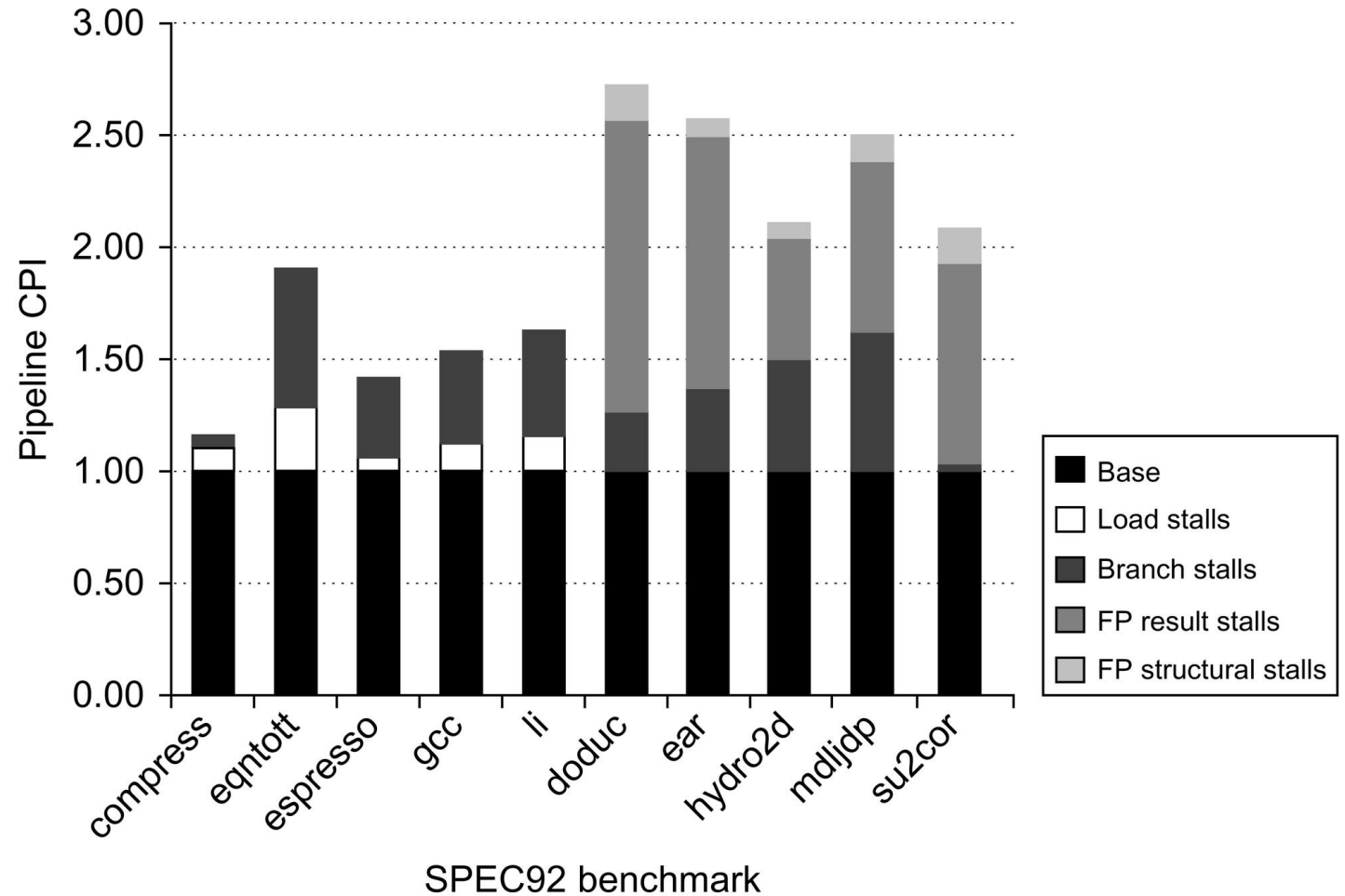
- **Data Dependences** : instructions are data dependent when there is a chain of RAW hazards between them.
- **Name Dependences** : instructions are name dependent when there is a WAR (anti-dependence) or WAW (output-dependence) hazard between them.

```
L.D    F0, 0(R1)
      ↓
ADD.D  F4, F0, F2
      ↙ ↘
L.D    F0, 0(R2)
```

- **Control Dependences** : Instructions dependent via branches.
if p1 { S1; }

R4000 Performance

- Non-Ideal CPI :
 - **Load stalls:**
1 ή 2 clock cycles
 - **Branch stalls:**
2 cycles + unfilled slots
 - **FP result stalls:**
RAW data hazard (latency)
 - **FP structural stalls:**
Not enough FP hardware (parallelism)



Instruction Level Parallelism (ILP)

- ILP: parallel execution of unrelated (independent) instructions
- gcc 17% control transfer instructions
 - 5 instructions + 1 branch
 - need to look beyond a code block to find more instruction level parallelism
- Loop level parallelism one opportunity
 - First SW, then HW approaches

FP Loop: where are the hazards?

```
while (R1 > 0) { M[R1] = M[R1] + F2; R1 -= 8 }  
Loop: L.D      F0,0(R1);F0=vector element  
      ADD.D    F4,F0,F2 ;add scalar from F2  
      S.D      F4,0(R1);store result  
      SUBI     R1,R1,8 ;decrement pointer 8B (DW)  
      BNEZ     R1,Loop ;branch R1!=zero  
      NOP                               ;branch delay slot
```

<i>Instruction producing result</i>	<i>Instruction using result</i>	<i>Latency in clock cycles</i>
FP ALU op	Another FP ALU op	3
FP ALU op	Store double	2
Load double	FP ALU op	1
Load double	Store double	1
Integer op	Integer op	0

Stalls?

Assume 5-stage RISC (in order)

FP Loop Showing Stalls

```

1 Loop: L.D   F0,0(R1)   ;F0=vector element
2         stall
3         ADD.D F4,F0,F2 ;add scalar in F2
4         stall
5         stall
6         S.D   F4,0(R1) ;store result
7         SUBI  R1,R1,8   ;decrement pointer 8B (DW)
8         BNEZ  R1,Loop  ;branch R1!=zero
9         stall          ;branch delay slot
  
```

<i>Instruction producing result</i>	<i>Instruction using result</i>	<i>Latency in clock cycles</i>
FP ALU op	Another FP ALU op	3
FP ALU op	Store double	2
Load double	FP ALU op	1

9 cycles per iteration:

Rewrite the code to minimize stalls!

Scheduled code for FP Loop

```
1 Loop: L.D    F0, 0(R1)
2          stall
3          ADD.D F4, F0, F2
4          SUBI  R1, R1, 8
5          BNEZ  R1, Loop    ;delayed branch
6          S.D   F4, 8(R1)   ;altered when move past SUBI
```

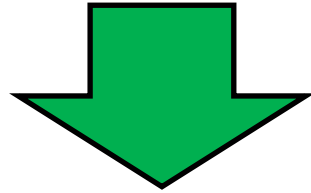
Move SD past BNEZ by modifying the address offset of SD

<i>Instruction producing result</i>	<i>Instruction using result</i>	<i>Latency in clock cycles</i>
FP ALU op	Another FP ALU op	3
FP ALU op	Store double	2
Load double	FP ALU op	1

6 cycles per iteration: Unroll loop 4 times to make code faster?

Loop Unrolling

```
while (R1 > 0) { M[R1] = M[R1] + F2; R1 -= 8 }
```



```
while (R1 >= 4*8) {  
  M[R1] = M[R1] + F2;  
  M[R1-8] = M[R1-8] + F2;  
  M[R1-16] = M[R1-16] + F2;  
  M[R1-24] = M[R1-24] + F2;  
  R1 -= 4*8  
}
```

```
while (R1 > 0) { M[R1] = M[R1] + F2; R1 -= 8 }
```

**Independent instructions
inside the loop. Good
opportunities for scheduling.**

Unroll Loop 4 times: name dependencies?

```
1 Loop: L.D    F0, 0(R1)
2      ADD.D   F4, F0, F2
3      S.D     F4, 0(R1)      ;drop SUBI & BNEZ
4      L.D     F0, -8(R1)
5      ADD.D   F4, F0, F2
6      S.D     F4, -8(R1)    ;drop SUBI & BNEZ
7      L.D     F0, -16(R1)
8      ADD.D   F4, F0, F2
9      S.D     F4, -16(R1)   ;drop SUBI & BNEZ
10     L.D     F0, -24(R1)
11     ADD.D   F4, F0, F2
12     S.D     F4, -24(R1)
13     SUBI    R1, R1, #32    ;alter to 4*8
14     BNEZ    R1, LOOP
15     NOP
```

Unroll Loop 4 times: name dependencies?

```
1 Loop: L.D    F0, 0(R1)
2      ADD.D   F4, F0, F2
3      S.D     F4, 0(R1)      ;drop SUBI & BNEZ
4      L.D     F0, -8(R1)
5      ADD.D   F4, F0, F2
6      S.D     F4, -8(R1)    ;drop SUBI & BNEZ
7      L.D     F0, -16(R1)
8      ADD.D   F4, F0, F2
9      S.D     F4, -16(R1)   ;drop SUBI & BNEZ
10     L.D     F0, -24(R1)
11     ADD.D   F4, F0, F2
12     S.D     F4, -24(R1)
13     SUBI    R1, R1, #32    ;alter to 4*8
14     BNEZ    R1, LOOP
15     NOP
```

How to deal with these?

No name dependencies now!

```
1 Loop: L.D    F0,0(R1)
2      ADD.D   F4,F0,F2
3      S.D     F4,0(R1)      ;drop SUBI & BNEZ
4      L.D     F6,-8(R1)
5      ADD.D   F8,F6,F2
6      S.D     F8,-8(R1)    ;drop SUBI & BNEZ
7      L.D     F10,-16(R1)
8      ADD.D   F12,F10,F2
9      S.D     F12,-16(R1)  ;drop SUBI & BNEZ
10     L.D     F14,-24(R1)
11     ADD.D   F16,F14,F2
12     S.D     F16,-24(R1)
13     SUBI    R1,R1,#32    ;alter to 4*8
14     BNEZ   R1,LOOP
15     NOP
```

“register renaming” removed WAR/WAW stalls

Unroll Loop 4 times

```
1 Loop: L.D    F0, 0(R1)
2      ADD.D   F4, F0, F2
3      S.D     F4, 0(R1)      ;drop SUBI & BNEZ
4      L.D     F6, -8(R1)
5      ADD.D   F8, F6, F2
6      S.D     F8, -8(R1)    ;drop SUBI & BNEZ
7      L.D     F10, -16(R1)
8      ADD.D   F12, F10, F2
9      S.D     F12, -16(R1)  ;drop SUBI & BNEZ
10     L.D     F14, -24(R1)
11     ADD.D   F16, F14, F2
12     S.D     F16, -24(R1)
13     SUBI    R1, R1, #32    ;alter to 4*8
14     BNEZ    R1, LOOP
15     NOP
```

1 cycle stall

2 cycles stall

eliminates overhead instructions,
but increases code size

Rewrite loop to
minimize stalls?

$15 + 4 \times (1+2) = 27$ cycles, or 6.8 cycles per iteration

Assumes R1 is multiple of 4

Schedule Unrolled Loop

```
1 Loop: L.D    F0, 0 (R1)
2         L.D    F6, -8 (R1)
3         L.D    F10, -16 (R1)
4         L.D    F14, -24 (R1)
5         ADD.D  F4, F0, F2
6         ADD.D  F8, F6, F2
7         ADD.D  F12, F10, F2
8         ADD.D  F16, F14, F2
9         S.D    F4, 0 (R1)
10        S.D    F8, -8 (R1)
11        S.D    F12, -16 (R1)
12        SUBI   R1, R1, #32
13        BNEZ   R1, LOOP
14        S.D    F16, 8 (R1) ; 8-32 = -24
```

14 cycles, or 3.5 cycles per iteration

- What kind of assumptions did we use to reorder and move the instructions?
 - OK to move store past SUBI even though changes register
 - **OK to move loads before stores/add: get right data?**
 - When is it safe for compiler to do such changes?

Compiler Perspectives on Code Movement

- Name Dependencies are hard to identify for Memory Accesses
 - $100(R4) == 20(R6)$?
 - for different iterations of the loop, is $20(R6) == 20(R6)$?
- In our example the compiler must understand that when R1 does not change then:

$$0 (R1) \neq -8 (R1) \neq -16 (R1) \neq -24 (R1)$$

- There were no dependencies between some loads and stores so they could be moved by each other

When is it safe to unroll a loop?

- **Example:** Are there dependencies? (A,B,C distinct & non-overlapping)

```
for (i=0; i<100; i=i+1) {  
    A[i+1] = A[i] + C[i];    /* S1 */  
    B[i+1] = B[i] + A[i+1]; /* S2 */  
}
```

- S2 uses the value, $A[i+1]$, computed by S1 in the same iteration.
- S1 uses a value computed by S1 in an earlier iteration, since iteration i computes $A[i+1]$ which is read in iteration $i+1$. The same is true of S2 for $B[i]$ and $B[i+1]$. This form of dependence (across iterations) is called **loop-carried dependence**
- In our prior example, each iteration was distinct. Dependences in the above example force successive iterations of this loop to execute in series.
- Implies that iterations can't be executed in parallel, right ?

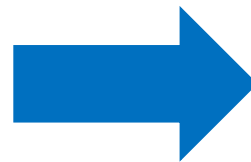
Loop-carried dependence: No parallelism?

- Example:

```
for (i=0; i<100; i=i+1) {  
    A[i]    = A[i] + B[i];    /* S1 */  
    B[i+1] = C[i] + D[i];    /* S2 */  
}
```

- S1 uses the value of B[i] which is produced by a previous iteration (loop-carried dependence).
- There is no other dependency. Hence, this dependence is not circular. We can conclude that the loop can be parallel.

A[0] = A[0] + B[0]
B[1] = C[0] + D[0]
A[1] = A[1] + B[1]
B[2] = C[1] + D[1]
A[2] = A[2] + B[2]
B[3] = C[2] + D[2]
...

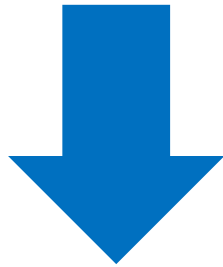


A[0] = A[0] + B[0]
B[1] = C[0] + D[0]
A[1] = A[1] + B[1]
B[2] = C[1] + D[1]
A[2] = A[2] + B[2]
B[3] = C[2] + D[2]
...

Loop-carried dependence: No parallelism?

- Example:

```
for (i=0; i<100; i=i+1) {  
    A[i] = A[i] + B[i];    /* S1 */  
    B[i+1] = C[i] + D[i]; /* S2 */  
}
```



```
A[0] = A[0] + B[0];    /* start-up code */  
for (i=0; i<99; i=i+1) {  
    B[i+1] = C[i] + D[i]; /* S2 */  
    A[i+1] = A[i+1] + B[i+1]; /* S1 */  
}  
B[100] = C[99] + D[99]; /* clean-up code */
```

Recurrence – Dependence Distance

- Example:

```
for (i=1; i < 100; i=i+1) {  
    Y[i] = Y[i-1] + Y[i];  
}
```

loop-carried dependence in **recurrence** form.

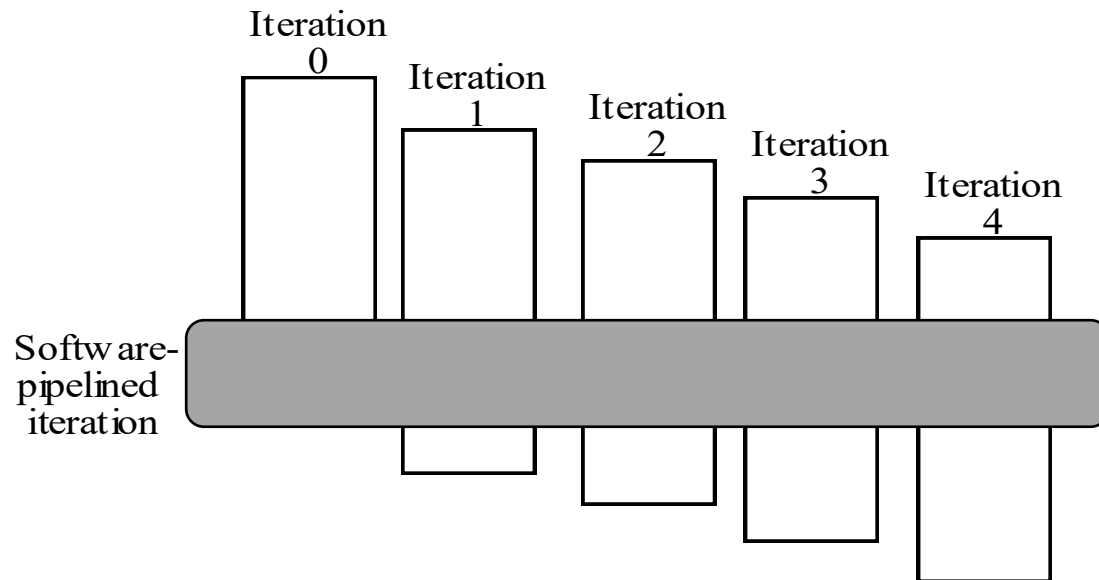
- Example:

```
for (i=5; i < 100; i=i+1) {  
    Y[i] = Y[i-5] + Y[i];  
}
```

Iteration i depends on iteration $i-5$, thus it has a **dependence distance** of 5. The longer the dependence distance the more potential to extract parallelism.

Alternative: Software Pipelining

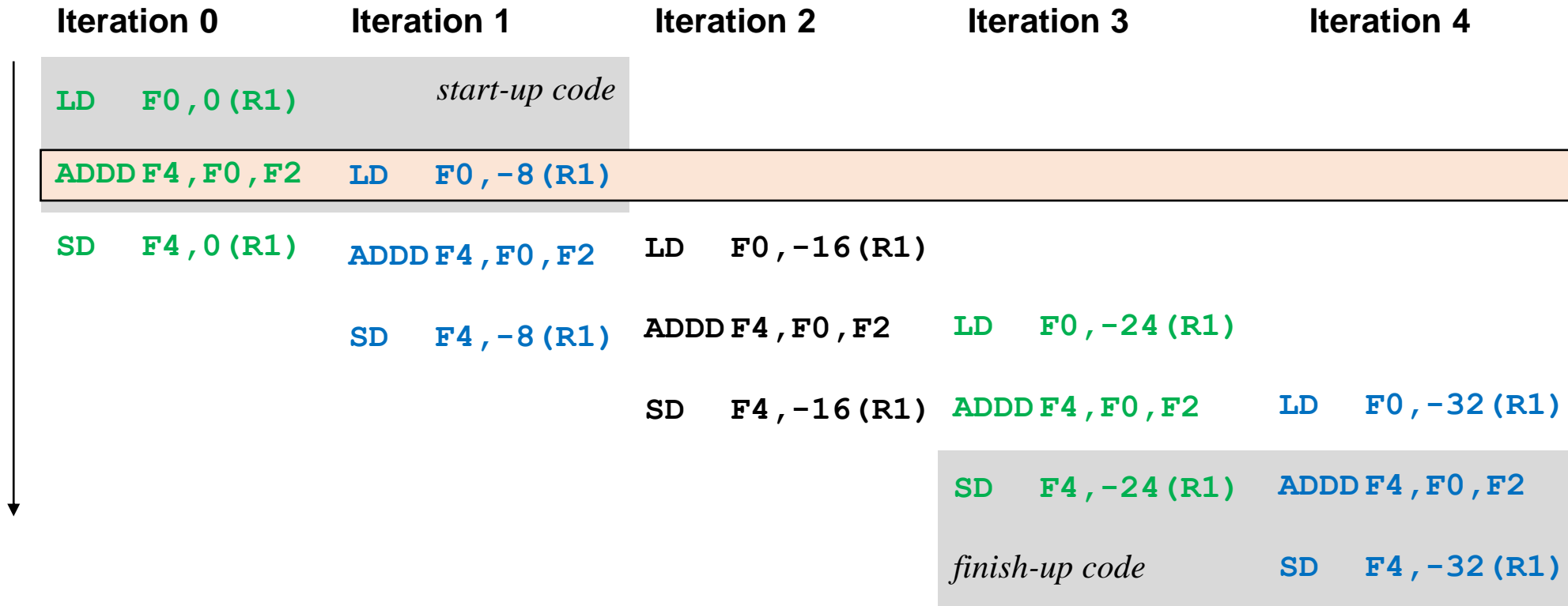
- Observation: If the iterations of the loop are independent, then we can exploit more ILP by executing instructions from different iterations of the loop.
- Software pipelining: reorganizes loops so that each iteration is made from instructions chosen from different iterations of the original loop without loop unrolling (Tomasulo in SW)



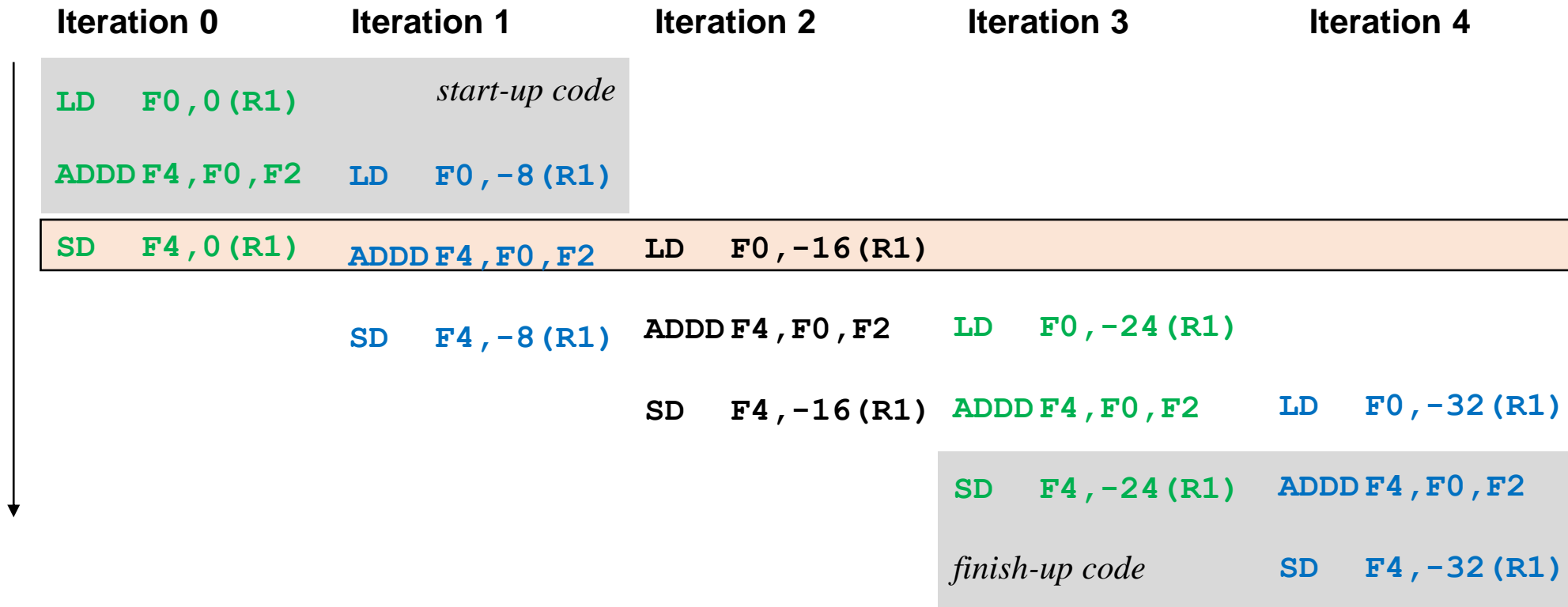
Software Pipelining: Example

Iteration 0	Iteration 1	Iteration 2	Iteration 3	Iteration 4
LD F0,0(R1)	<i>start-up code</i>			
ADDD F4,F0,F2	LD F0,-8(R1)			
SD F4,0(R1)	ADDD F4,F0,F2	LD F0,-16(R1)		
	SD F4,-8(R1)	ADDD F4,F0,F2	LD F0,-24(R1)	
		SD F4,-16(R1)	ADDD F4,F0,F2	LD F0,-32(R1)
			SD F4,-24(R1)	ADDD F4,F0,F2
			<i>finish-up code</i>	SD F4,-32(R1)

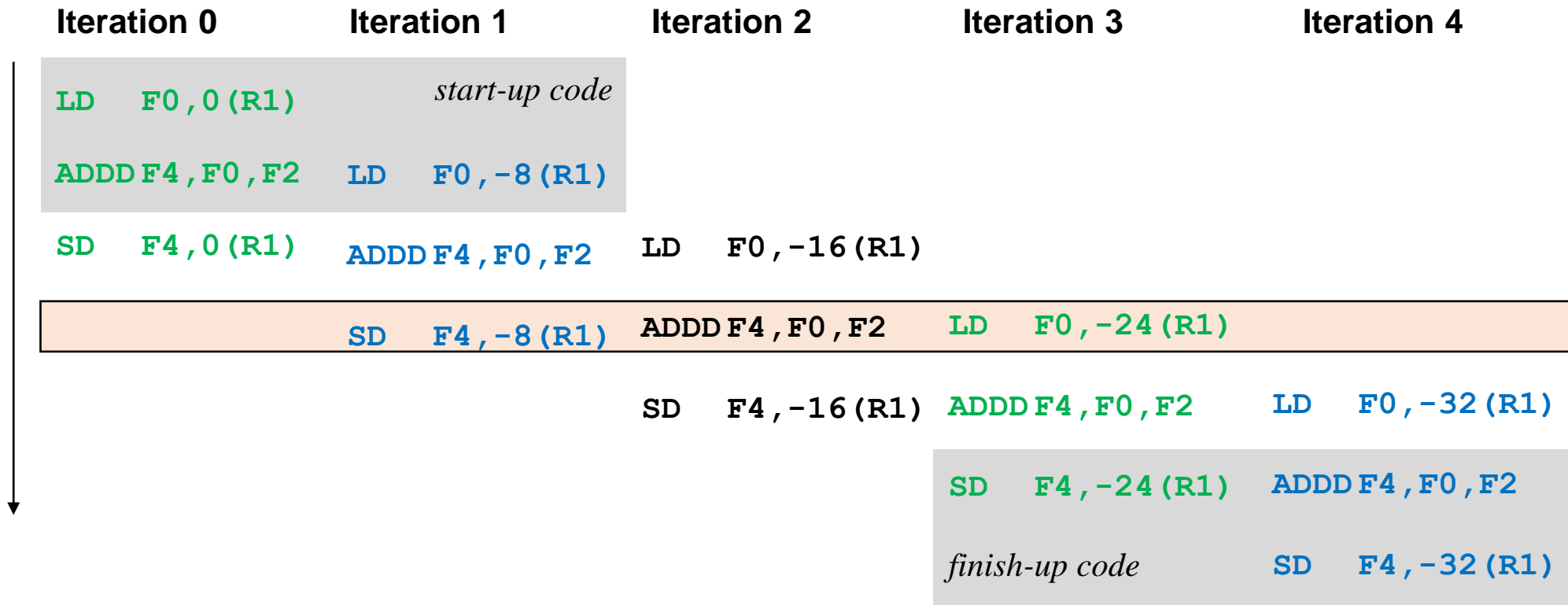
Software Pipelining: Example



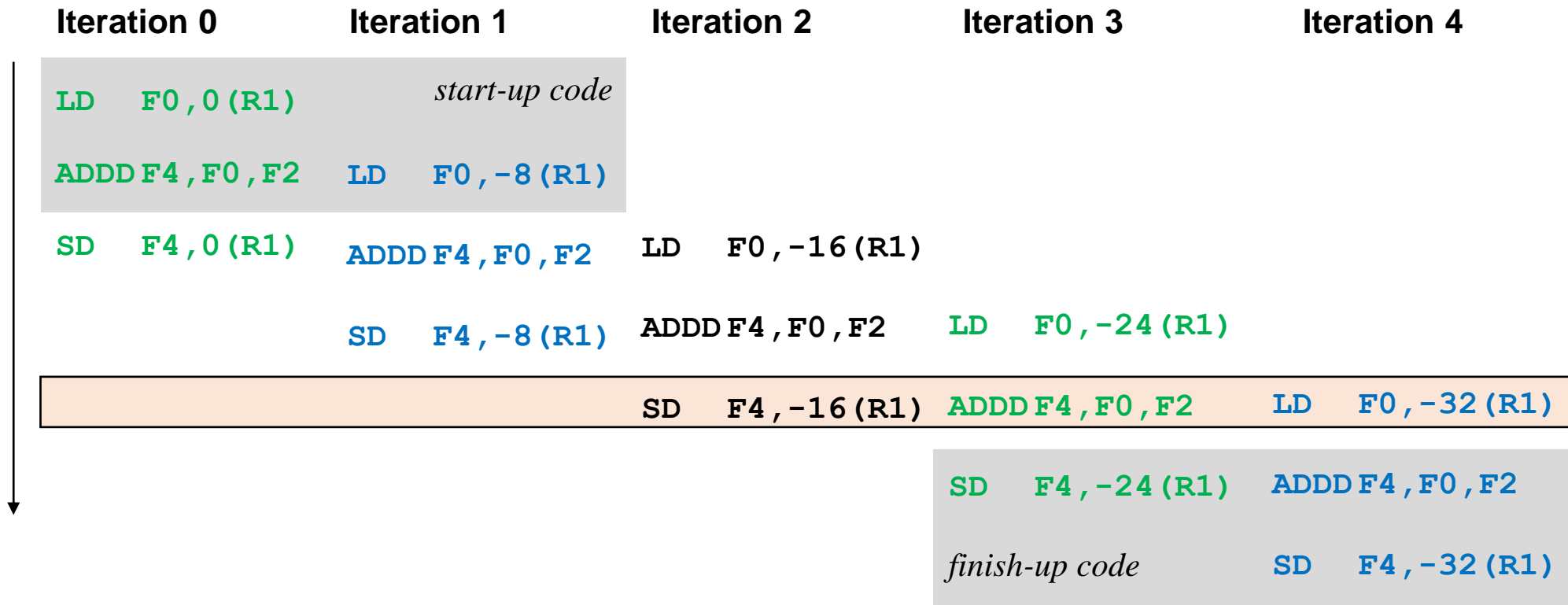
Software Pipelining: Example



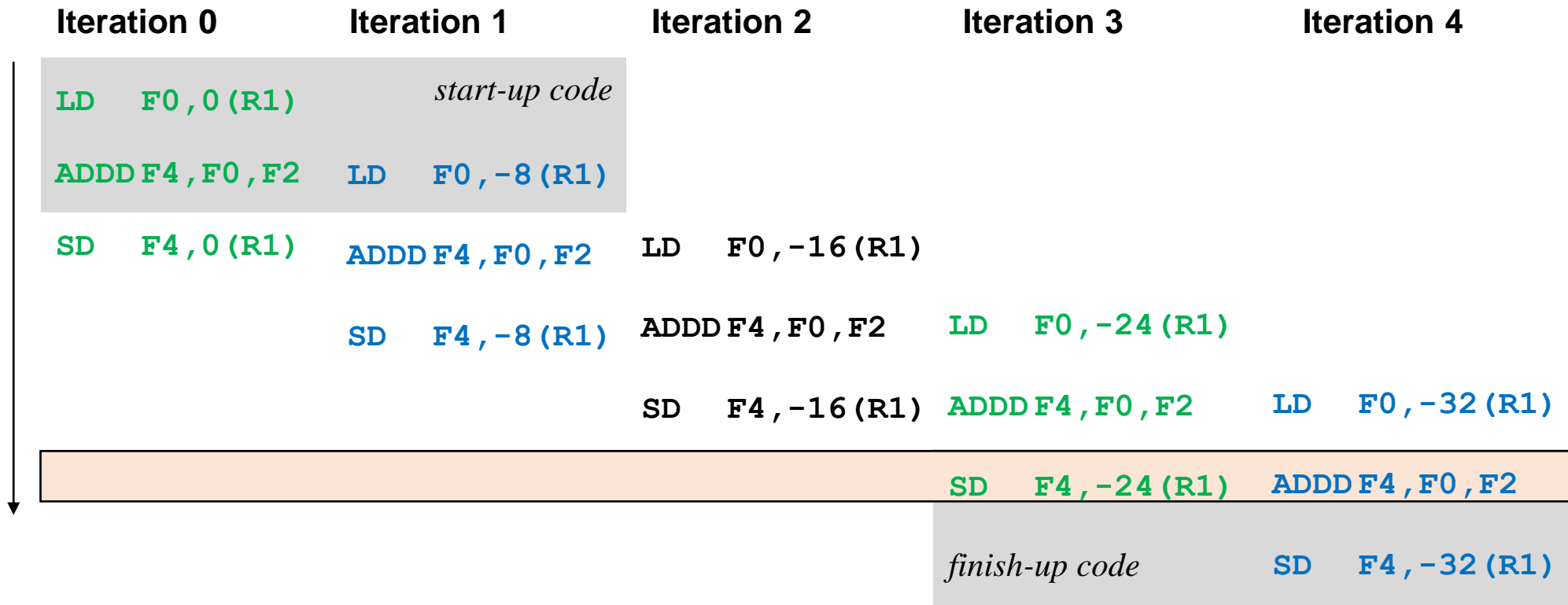
Software Pipelining: Example



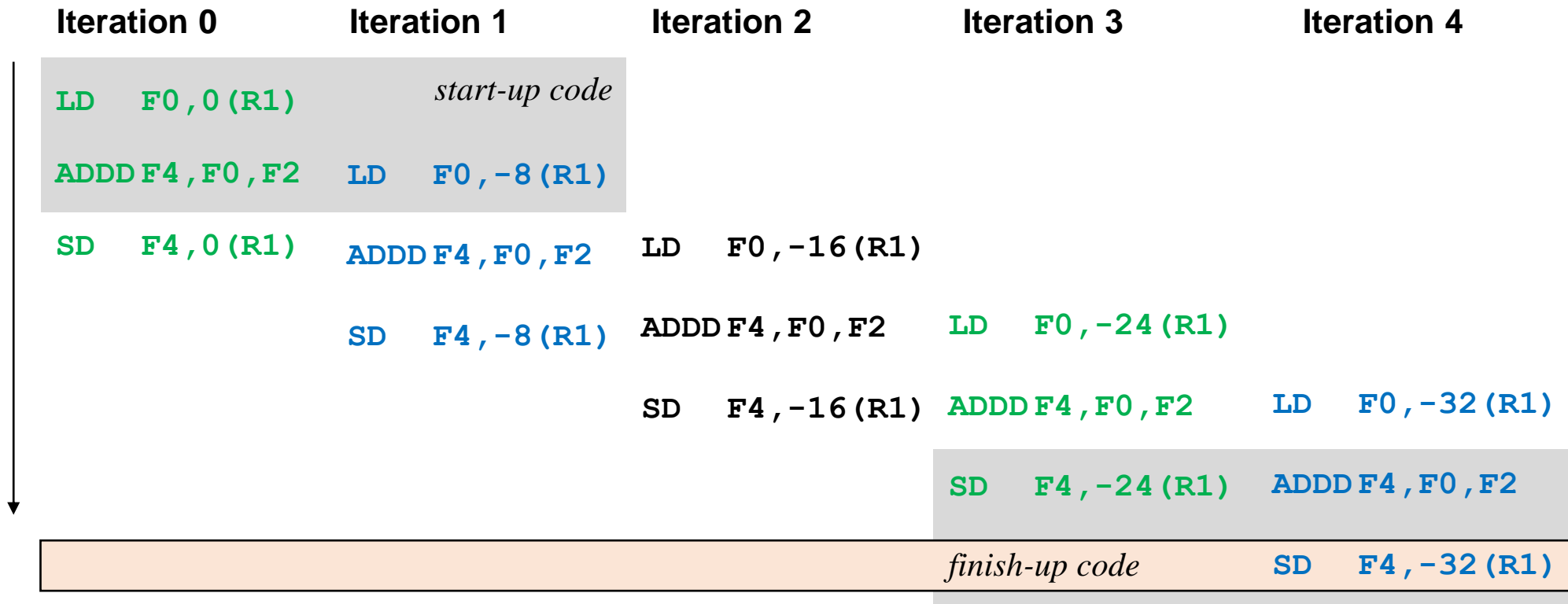
Software Pipelining: Example



Software Pipelining: Example



Software Pipelining: Example



Software Pipelining: Example

Before: Unroll 3 times

```
1  L.D   F0,0(R1)
2  ADD.D F4,F0,F2
3  S.D   F4,0(R1)
4  L.D   F6,-8(R1)
5  ADD.D F8,F6,F2
6  S.D   F8,-8(R1)
7  L.D   F10,-16(R1)
8  ADD.D F12,F10,F2
9  S.D   F12,-16(R1)
10 SUBI  R1,R1,#24
11 BNEZ  R1,LOOP
```

After: Software Pipelined

```
1  S.D   F4,0(R1) ; Stores M[i]
2  ADD.D F4,F0,F2 ; Adds to M[i-1]
3  L.D   F0,-16(R1) ; Loads M[i-2]
4  SUBI  R1,R1,#8 ; i = i - 1
5  BNEZ  R1,LOOP
```

5 cycles per iteration

RAW hazards convert to WAR hazards.

Software Pipelining vs Loop Unrolling

Symbolic Loop Unrolling

- **Maximize result-use distance**
- **Less code space than unrolling**

But..

- **Harder to implement**
- **Execution of SUB & BNEZ in every iteration**